

# Key Generation Based on Henon Map and Lorenz System

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## ABSTRACT

Securing information has been the most significant process for communication and data store. Orderly to secure information such as data authentication, data integrity, and confidentiality must be verified based on algorithms of cryptography. Where, the most important part of any encryption algorithms is the key which specifies if the system is strong enough or not. The proposal of this paper is a new method to generate keys based on two kinds of chaos theory in order to improve the security of cryptographic algorithms. The base of this proposal is to investigate a new method for generating random numbers by using the 3D Lorenz system and 2D Henon map. The newly generated keys have successfully passed the National Institute of Standards and Technology (NIST) statistical test suite.

**KEYWORDS:** 2D Henon map; 3D Lorenz system; NIST test suite.

## الخلاصة

يعتبر تأمين المعلومات العملية الأكثر أهمية لغرض اتمام التواصل وتخزين المعلومات. من اجل تأمين المعلومات كالتحقق من صحة البيانات وتكامل البيانات وسريتها يتم استخدام خوارزميات التشفير الجزء الأكثر أهمية في أي خوارزمية تشفير هو المفتاح والذي يحدد ما اذا كان النظام قويا كفاية ام لا. نقترح في هذه الورقة طريقة جديدة لإنشاء المفاتيح بالاعتماد على نوعين من نظريات الفوضى من اجل تحسين امن خوارزميات التشفير. اساس هذا الاقتراح هو اكتشاف طريقة جديدة لإنشاء ارقام عشوائية باستخدام خريطة لورنز ثلاثية الابعاد وخريطة هينون ثنائية الابعاد. اجتازت المفاتيح التي تم انشاؤها حديثا بنجاح مجموعة الاختبارات الاحصائية للمعهد الوطني للمعايير والتقنية.

## INTRODUCTION

The quantum of information exchanged through the network was increased with the development of communication networks, and the dependence of organizations on these new communication channels has grown dynamically. At the same time, the risks are incremented significantly, so, many technologies like cryptography were developed to overcome these threats[1]. The cryptography uses to secure data, which means a mechanism to protect the data over communication network [2]. In cryptography, there are two kinds of cryptographic algorithms: symmetric and asymmetric cryptography [3]. In symmetric cryptography, the same key is shared between the sender and receiver e.g. Data Encryption Standard (DES) and Advanced Encryption Standards (AES)[4]. While in asymmetric cryptography sender and receiver used different keys to encrypt and decrypt data [3]. the first key, which is called the private key, is kept secret and another one known as the public key

is revealed and this removes the need for the sender and the receiver to share key. The only request is that public keys are shared with the users who are authenticated [5]. RSA (Ron Rivest, Adi Shamir, and Leonard Adleman) which is considered as a base to the e-commerce revolution [6], is one of the public key algorithms which is widely used in this context [4].

For a lot of reasons, all these an encryption algorithm typically uses a pseudo-random encryption key generated by a key generation process[7]. Chaotic signal has been much interesting in the past few decades for used it vastly in cryptography. The popularity of chaotic systems depends on the non-predictable behavior and randomness of these systems [8]. The most important feature of the chaotic system is sensitivity to the initial condition; Edward Lorenz is the first researcher who clarifies this property. During a search, he finds out that if the initial conditions in the system of differential equations bit changed that would

completely after the short time change the resulting, and many researchers have been confirming this special feature of chaos[9]. In this paper, a new key generator will be presented based on the 2D Henon map and 3D Lorenz system.

The paper topics can be illustrated as shown below: the chaotic map was briefly introduced in section II. the related work implemented in section III, In the following two sections proposed work, results and discussions were explained. and finally, the conclusions.

## CHAOSE SYSTEMS

Chaos is irregular long-term behavior in a deterministic system that exhibits sensible dependence on initial states[10].

### Henon Map

the famous two-dimensional Henon map was proposed by H'enon in 1978, as a diminutive approach to study the dynamics of the Lorenz system[11]. The study of the Henon map shows a simple two-dimensional map with quadratic nonlinearity equation. The map gives a first example of the exotic attractor with a fractal structure[12]. which is described as following:

$$x = 1 - ax^2 + by \quad (1)$$

$$y = x \quad (2)$$

Here, the dynamic behavior of the system depends on the values of the parameters, a and b. The system is chaotic when the value of a and b are 1.4 and 0.3 as shown by Figure 1, while x and y are variables [8].

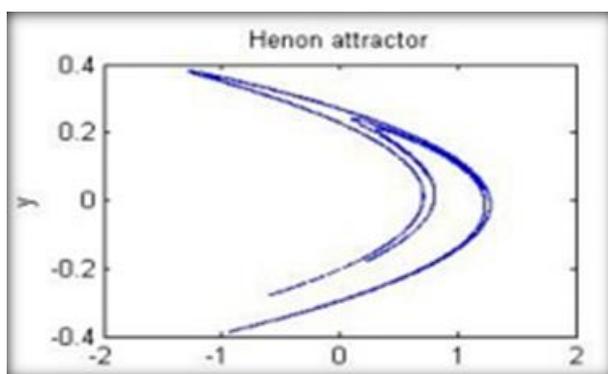


Figure1. Chaotic behavior of Henon map.

### Lorenz System

About 1960, Edward Lorenz invented the Lorenz system which is a dynamical system

characterized by a nonlinear system of, simple differential equations[14].

$$x' = \sigma(y - x) \quad (3)$$

$$y' = rx - y - xz \quad (4)$$

$$z' = xy - bz \quad (5)$$

$\sigma$ ,  $b$  and  $r$  are parameters. the system enters a chaotic scope when Choosing  $\sigma = 10$ ,  $r = 28$ ,  $b = 8/3$ . So, given initial values  $x_0$ ,  $y_0$ , and  $z_0$ , the system will speedily spread and generate values widely different from a system given only little different values for  $x_0$ ,  $y_0$ , or  $z_0$  [13].

## RELATED WORK

In 2018, Ali Kashmar proposed a new method for key stream generator Based on Chebyshev Maps to meet the requirement of image encryption. The output is tested in different measurements; outcomes show that our proposed method is resistant versus security analysis and has good cryptographic strength[9].

In 2017, Ronald Marsh, Scott Kerlin proposed a 'many-key approach' for images using the Lorenz System that combines bit by bit diagonal and anti-diagonal mix of the pixels in an image, using a large key space of 2808 bits for 24-bit color images and 936 bits for gray scale images. However, as shown, a large key length is not sufficient .some form of diffusion can be used to provide high security which provides very secure image encryption [13].

In 2017, Ekhlas Abbas Albahrani, Tayseer Karam proposed a new schema based on a combination of two types of chaotic maps which is 3D Cat map and 3D Henon map to generate keystream .the first step in this method is using 3D Henon map to generating random numbers and converted these numbers to a binary sequence. Then, in the final step permuted and XOR the generated sequence positions by using the 3D Cat map. The outcome of this schema is the goodness of the produced keystream, and a high degree of security versus different attacks, sensibility to the initial values [15].

In 2016, G. Madhuri, I.M.V. Krishna proposed schema to provide efficient and secure key generation uses a non-linear chaos theory approach. The secret keys K1 and K2 have many potential choices that give the image the highest level of security. the master key is with a length of 64 bits binary string. The two-session keys

generated from the master key are each 64 bit long, and are processed against versus the input image turned into bit blocks of 64 bits each[16]. In 2013, N. S. Raghava, Ashish Kumar Henon's map with byte sequences and a new approach of pixel shuffling was used to suggest a new symmetric encryption algorithm for image encryption. the suggested method resulting in effective encryption of images[8].

### Proposed Key Generation System

In the proposed system, two of chaotic maps are used 2D Henon map and 3D Lorenz system, it will generate a sequence of random numbers between zero and one, these numbers will be processed according to a specific mechanism and generate keys as illustrate in the steps bellow. Figure 2 indicates the generation system method.

**Step1:** Apply 2D Henon map and 3D Lorenz chaotic system.

For i=1 to 1000

$$\left. \begin{aligned} x &= 1 - ax^2 + y \\ y &= bx \end{aligned} \right\} \text{Henon map}$$

$$\left. \begin{aligned} x' &= \sigma(y - x) \\ y' &= rx - y - xz \\ z' &= xy - bz \end{aligned} \right\} \text{Lorenz system}$$

**Step 2:** Set the numbers produced by x in string 1, y in string 2, x' in string 3, y' in string 4, z' in string 5.

**Step 3:** Remove the negative sign.

**Step 4:** Cut number s after the comma.

**Step 5:** Convert number to hexadecimal

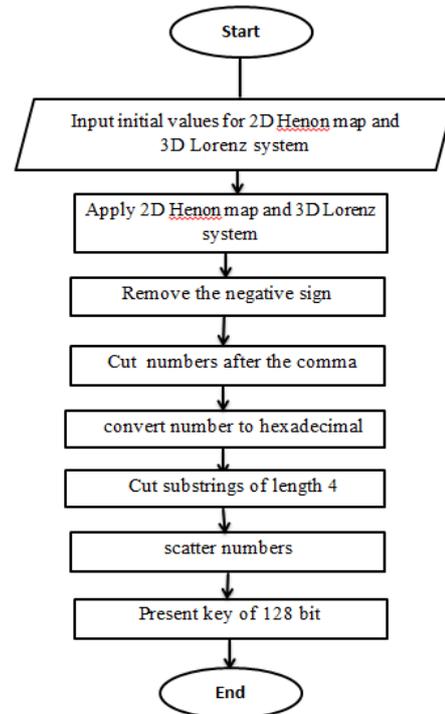
**Step 6:** cut substrings of length 4 from each string.

**Step 7:** scatter numbers by taking the first two digits from each string then the second two digits.

**Step 8:** divide the result of step6 to sub string of length 32

**Step 9:** convert result of step7 to binary of length 4.

**Step 10:** Present key of 128 bit.

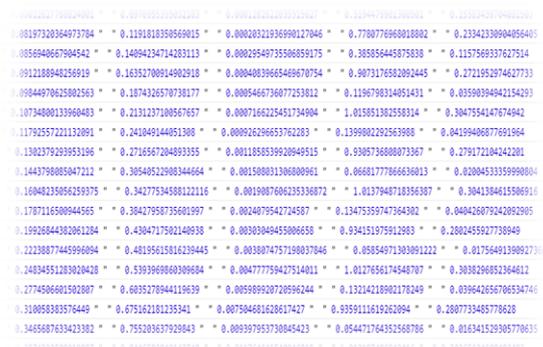


**Figure2.** Block diagram illustrate Steps of key generation.

**Step1:** Apply 2D Henon map and 3D Lorenz chaotic system.

Input the initial values and the control parameters values which are a =1.76 and b =0.1 to the Henon map equations and  $\sigma =10$ ,  $r=28$ ,  $b =8/3$  to the Lorenz system equations. The initial values for the Henon map  $(x_0,y_0)$  and Lorenz system  $(x_0,y_0,z_0)$  are floating-point numbers with a range of [0-1].

The output of this step is floating numbers. In **step2** set, the numbers produced by x in string 1, y in string 2, x' in string 3, y' in string 4, z' in string 5. and **step 3** remove the negative sign, Figure 3 represents a sample of 20 iterations, the first three columns are for the Lorenz system and the second two-column is for the Henon map.



**Figure3.** The output of Henon map and Lorenz system.

**Step 4: Cut number s after the comma.**

Cut numbers after the comma and take only float numbers, as shown in Figure 4 then convert numbers to hexadecimal in step 5 while, In step 6 take only substrings of length 4 from each string to obtain high randomness, as shown in Figure 4 and Figure 5.

```

12627798624001 978959335032183 1282622653515627 5194479901360301 13583439704001504
197320364973784 1191818350569015 20321936990127050 7780776968018802 23342330904056404
8940667904542 14094234714283112 29549735506859176 3858564545875838 1157569337627514
2188948256919 16352700914902918 40839665469670750 9073176582092444 2721952974627733
44970625802564 1874326570738177 5466736077253812 1196798314051431 3590394942154293
734800133960484 2131237100567657 7166225451734904 15851382558314 3047554147674942
792557221132092 241049144051308 926296653762283 1399802292563988 4199406877691964
302379293953196 2716567204893355 11858539920949516 9305736808073368 279172104242201
443798085047212 30540522908344664 1508031306800961 668177866636013 200453359990804
8048235056259376 34277534588122116 19087606233536870 137948718356387 3041384615506916
787116500944565 38427958735602000 20079542724587 13475359747364302 40426079242092904
392644382061284 4304717502140938 303049455006658 934151975912983 2802455927738949
2238877445996096 48195615816239450 38074757198037850 5854971303091222 17564913909273664
4834551283020428 5393969860309684 4777759427514011 127656174548707 3038296852364612
774506601502807 6035278944119639 5989920720596244 13214218902178248 39642656706534744
10058383576449 675162181235341 7504681628617427 9359111619262094 2807733485778628
465687633423382 755203637929843 9397953730845424 54471764352568780 16341529305770636
    
```

Figure4. The float numbers.

```

7d9c" "12ea" "10f3" "9b74" "4a0b"
8f51" "14e3" "1a77" "1543" "3fca"
a268" "1725" "a187" "1816" "739f"
7280" "19b5" "9de2" "13e3" "3baa"
ce04" "2dc6" "9a80" "19e8" "4db8"
e731" "1fdf" "976b" "a882" "328d"
1031" "2b8d" "94ad" "1d64" "582d"
1220" "41f0" "ea21" "a4cd" "1ee6"
1447" "8b9b" "e71f" "226f" "674f"
16ad" "de0e" "e500" "574a" "29e6"
1959" "1f6d" "8e76" "233c" "a921"
1c54" "1a16" "8e99" "1c43" "54ca"
1fa8" "214a" "8fb6" "225d" "107e"
235f" "5cf8" "9203" "6fe2" "14fa"
3ff3" "17ae" "17f6" "234c" "69e4"
8a14" "18d7" "9b54" "1dd2" "8f24"
1616" "872b" "104f" "2248" "66d8"
1394" "14de" "115c" "502b" "f083"
    
```

Figure5. The four digits random numbers.

**Step 7: Scatter Numbers**

This step is worked for scattered the numbers to arrive high randomness for key generated by cutting the first two digits from x0, y0, z0 (Lorenz system) and x0, y0 (Henon map) save the result in array. Then, considering the second two digits are proposed and so on until the end of iterations, as shown in Figure 6.

**Step 8: divide the result of step6 to sub string of length 32**

Each substring then will convert to binary in step 9.

**Performance and Security Analysis**

In order to analyze the performance and measure the degree of security of our proposal, some

cryptographic tests must be applied such as randomness tests.

```

3f101210e5471111209b7d12109b488f141a153fa217a1187372199a133bce2d9a194de71f97a832102b941d581241eaa41
661314150f0292b1223a921b4201f975b27162266ebd2191b6272f1d5a11962e222164cc175022661a52c923ab0a782382
16112e4433ab32542028792d361621a141037611d87327341ab6b2246437d31142e126412b31b74232436574c1b3348202
ac183b3d2c8663211621a011385aa03024202d61ac394e462d81b323221a00ab365f10312b121b63acca3a3248d92033482
da201c2c1648f703211332a2eeea66ac1414372d0a8ed43621a02238411133104e2466ac362e3048da2637202164254041a
a01f14fac33e62d30a4acae12192ddad7ef5735a01a1d271133202d14a4ac1858342d0881d172a2164382a2b11339a2d2ca
33201e8c666b15b1b2d0884e318535a0211d761b3332d924666b1729352d8835d53b21641a252e1b5338592666b1b0862a4
ac2349252d083268282164dafb2811331a127e66ac24c3132dda2e122d21643823611339e173c66ac233c684da2c2c2
15f0ac7e21a05965511331fc613666b13122c2dda40b6cb21a0a0b62711338dbd96666b11c93c2d0814db992164174271
64261f1dac331c6529666b2114192dda2a16a621a08220351153e2152766ac34211b2dda1b61121a023157f1133064b3b6
331d3524a46b38381d48da26240216417414911331054e0666b373632d152b42442164201d2e1b33133d6666b45291b7
11b0281348da2e1c4a21a019a82611334ee1d1a4ac2a12c48da1a1f24216418325611338da4e3266ac77f72c48da7a22c3
882123021641e001e1b534810e666b1119b5408da1d24221a02930371133534c1966ac8d121348da15221721102474b1d
a02713381133362d2166ac7238c2dda1a262a2164263e141133322359a4ac322c22d0881211022164112a2f331528166
    
```

Figure6. Scattered the numbers.

**Statistical Analysis**

The output should offer a high degree of randomness, whatever the initial values. So, statistical analysis should be assortment to show the nature and quality of the binary sequences.

**Randomness test**

Randomness test is done on binary sequences and applied through statistical tests suite NIST [8]. The result from step 3 was converted to binary sequences. 1000 of various binary sequences each of which has length 128bits was passed through the NIST test. NIST consists of 5 tests which are frequency-test, run-test, poker-test, serial-test, auto-correlation test. Each test has a p-value which is compared to fixed significance level  $\alpha$ . If the p-value  $\leq \alpha$ , the test is passed otherwise it is failing. The result of a statistical test is illustrated in Table 1.

Table 1: Five test's result of(128 bits) key space

Tests		Freedom Degree	State
Frequency Test		$\leq 3.84$	Pass =1.531
Run Test	T0	$\leq 13.784$	Pass =12.938
	T1	$\leq 7.531$	Pass = 4.188
Poker Test		$\leq 11.1$	Pass =4.725
Serial Test		$\leq 7.81$	Pass =2.875
Auto Correlation Test	1 <sup>st</sup> Shift	$\leq 3.84$	Pass= 0.638
	2 <sup>nd</sup> Shift		Pass= 3.175
	3 <sup>rd</sup> Shift		Pass= 0.392
	4 <sup>th</sup> Shift		Pass=0.806
	5 <sup>th</sup> Shift		Pass=0.984
	6 <sup>th</sup> Shift		Pass=1.180
	7 <sup>th</sup> Shift		Pass=0.008
	8 <sup>th</sup> Shift		Pass=0.133
	9 <sup>th</sup> Shift		Pass=0.008
	10 <sup>th</sup> Shift		pass=0.034

**Key space analysis**

It's simple: if you want to avoid brute-force attacks and want cipher to be secure, then the proposed system should have enough possible keys (a large enough key space) that an attacker cannot simply try every one of them. Cryptographers are conservative, so they usually specify key space to be more than  $2^{128}$  possible keys (128-bit security). In the proposed system the key space is consist of the Lorenz system parameters(x0, y0, z0) and Henon map parameters(x0, y0). If each of these parameters has a precision of  $10^{-16}$ , the key space size for the Lorenz system will be  $2^{192} ((10^{16})^3)$  for initial values. And the key space size for initial values to the Henon map is  $2^{128} ((10^{16})^2)$ . Finally, the full space of keys is  $2^{192} + 2^{128} = 2^{320}$ .

**Key sensitivity analyses**

To warranty the sensitivity of keys Correlation test is used which is to check the correlation between the generated sequences of keys. Correlation test involves two ways:

**The Pearson's correlation coefficient** which means calculating the correlation coefficient among each pair of created keys sequences to analyzes the correlation between them. In a sample, it is denoted by  $R_{(str, str1)}$  and Hence, The formula described below is used to find the Pearson R correlation [15]:

$$R_{(str, str1)} = \frac{\sum_{i=0}^{n-1} (x_i - \bar{x}) \cdot (y_i - \bar{y})}{[\sum_{i=0}^{n-1} (x_i - \bar{x})^2]^{1/2} \cdot [\sum_{i=0}^{n-1} (y_i - \bar{y})^2]^{1/2}} \quad (6)$$

str, str1 represent two sequences specified by  $Str = [x_1, \dots, x_n]$  and  $Str1 = [y_1, \dots, y_n]$ .

Where:

$\bar{x} = \sum_{i=0}^{n-1} x_i / n$ ,  $\bar{y} = \sum_{i=0}^{n-1} y_i / n$  represents the mean values of Str and Str1.

**Hamming distance** is the second method of correlation test to warranty the sensitivity of keys, this way of correlation based on analyzing directly the bits of sequences.by finding the number of different bits between two binary sequences that have the same length M. Thus, for two binary sequences, the Hamming distance is given by:

$$d(s_1^b, s_2^b) = \sum_{j=0}^{m-1} (x_j \oplus y_j) \quad (7)$$

where  $x_j$ ,  $y_j$  are the elements of  $S_1^b$ ,  $S_2^b$ . The binary sequences are truly random when the

normal distance is about M/2, who gives an attribution of around 0.50 [17].

With the pseudo-random generator, the sensitivity of keys is an essential element. That's mean, the output must be uncorrelated if a little different happened on initial values. Correlation tests (Pearson's correlation and Hamming distance) are performed to warranty the sensitivity of the key on four binary key sequences K1, K2, K3, K4, which are generated from a little various in initial values. Table 2 illustrates the results of Pearson's correlation coefficients and Hamming distance which explain that the correlation between the generated sequences is little.

**Table 2:** The results of Pearson's correlation coefficients and Hamming distance.

	Tests	
	Correlation Analysis	Correlation Analysis
<b>K1vs k2</b>	0.316395	0.316395
<b>K1vs k3</b>	0.031362	0.031362
<b>K1vs k4</b>	-0.24205	-0.24205
<b>K2vs k3</b>	0.110193	0.110193
<b>K2vs k4</b>	0.317937	0.317937
<b>K3vs k4</b>	0.150229	0.4765625

**Speed analysis**

The analysis of speed performance is done on a personal computer which characterized by Intel(R) Core(TM) i7 CPU M620@ 2.67 GHz 2.67 GHz. The algorithm is implemented using JavaScript on the Komodo editor. In this analysis, binary sequences of 128-bit lengths are generated and its execution time is calculated. Table 3 shows the performance time in milliseconds.

**Table 3:** Performance analysis.

Length in Bit	Speed in milliseconds
128 bit	631

**CONCLUSIONS**

A chaotic keys generator was proposed in this paper based on the 2D Henon map and 3D Lorenz system. The initial conditions (x0,y0) are the input to the 2D Henon chaotic map and (x0, y0, z0) are the input to the 3D Lorenz chaotic system. float numbers are only taken and used only four digits from it. Scattered is used to obtain high randomness by taking the first two digits from these four digits and then the second two digits until the end of numbers.

The proposed system has the ability to generate a large number of key sequences that can be useful in many applications in cryptography. The system has the sensitivity to the initial values (keys), the quality of the produced key sequences and the degree of security versus several attacks.

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